

Math 776 Graph Theory Lecture Notes

Extremal Graph Theory

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Theorem 1 (Mantel 1907) *The maximum number of edges in an n -vertex triangle-free simple graph is $\lfloor n^2/4 \rfloor$.*

Definition 1 *A graph G is **H-free** if G has no subgraph isomorphic to H .*

A triangle-free graph is C_3 free.

Example 1 *Bipartite graphs are C_3 free since they contain no odd cycles.*

Proof of Theorem 1: Let G be a triangle-free graph. Let x be a vertex of G with maximum degree. Let $k = d(x)$ be the degree of x . Let $N(x)$ be the set of neighbors of x . Since G is triangle-free, there are no edges whose endpoints are both in $N(x)$. $\overline{N(x)}$ forms a vertex cover. The number of edges $e(G)$ satisfies

$$\begin{aligned} e(G) &\leq \sum_{y \in N(x)} d(y) \\ &\leq \sum_{y \in N(x)} k \\ &= k \cdot |N(x)| \\ &= k(n - k) \\ &\leq \left(\frac{k + (n - k)}{2} \right)^2 \\ &= \frac{n^2}{4} \end{aligned}$$

Since $e(G)$ is an integer,

$$e(G) \leq \lfloor \frac{n^2}{4} \rfloor.$$

This upper bound is reachable by

$$K_{\lfloor \frac{n}{2} \rfloor, \lceil \frac{n}{2} \rceil}.$$

$$\lfloor \frac{n}{2} \rfloor \lceil \frac{n}{2} \rceil = \lfloor \frac{n^2}{4} \rfloor$$

Case 1: n is even

$$n = 2k$$

So,

$$\lfloor \frac{n}{2} \rfloor \lceil \frac{n}{2} \rceil = k^2$$

and

$$\lfloor \frac{n^2}{4} \rfloor = k^2$$

Case 2: n is odd

$$n = 2k + 1$$

So,

$$\lfloor \frac{n}{2} \rfloor \lceil \frac{n}{2} \rceil = k(k + 1) = k^2 + k$$

and

$$\lfloor \frac{n^2}{4} \rfloor = \lfloor \frac{(2k + 1)^2}{4} \rfloor = \lfloor k^2 + k + \frac{1}{4} \rfloor = k^2 + k$$

□

Definition 2 A **Turán Graph** $T_{n,r}$ is a complete r -partite graph with n vertices whose partite sets differ in size by at most 1.

Theorem 2 (Turán, 1941) Among the n -vertex simple graphs with no $r + 1$ clique, $T_{n,r}$ has the maximum number of edges.

Lemma 1 Among simple r -partite graphs with n -vertices, the Turán graph is the unique graph with the most edges.

Proof of Lemma: Let G be a simple r -partite graph with the most edges.

1. G must be a complete r -partite graph.
2. Any two partite sets of G differ in size by at most one.

Prove 2. by contradiction. Say $|v_1| \leq |v_2| - 2$. Pick any vertex in v_2 and move it up into v_1 . Call this graph G' .

$$\begin{aligned} e(G') &= e(G) - dx^{old} + dx^{new} \\ &= e(G) - (n - |v_2|) + (n - |v_1| - 1) \\ &= e(G) + |v_2| - |v_1| - 1 \\ &> e(G) \end{aligned}$$

Contradiction.

Say $n = ar + b$, $0 \leq b \leq r$. So, $n = \underbrace{(a + 1) + \dots + (a + 1)}_b + \underbrace{a + \dots + a}_{r-b}$.

Hence, $T_{n,r}$ has the most edges among r -partite simple graphs. □

Proof of Turán Theorem:

Claim: Among graphs with no $r + 1$ clique, the maximum is achieved by an r -partite graph. For each G , construct H satisfies:

1. H is r -partite.
2. $e(G) \leq e(H)$.

Prove by induction on r .

Initial Case:

($r=1$): G has no edges, i.e. G is K_2 free. This case is trivial.

($r=2$): Mantel's Theorem

Inductive Hypothesis:

Suppose the claim is true for r . For $r+1$, let G be a graph containing no K_{r+2} . Let x be a vertex with maximum degree k . The induced subgraph on $N(x)$ contains no K_{r+1} . By inductive hypothesis, there is an r -partite graph H' with $e(G_{N(x)}) \leq e(H')$. Let H be the graph joining $\overline{N(x)}$ and H' .

1. H is an $r+1$ partite graph.
- 2.

$$\begin{aligned}
 e(H) &= |\overline{N(x)}|k + e(H') \\
 &\geq \sum_{y \in N(x)} dy + e(G_{N(x)}) \\
 &\geq e(G).
 \end{aligned}$$

By the lemma, $T_{n,r}$ has the most edges. □

Definition 3 Turán number $t(n, G)$ is the maximum number of edges that an n -vertex G -free simple graph can have.

$$\begin{aligned}
 t(n, K_r) &= e(T_{n,r-1}) \\
 &= \binom{n}{2} - b \binom{a+1}{2} - (r-1-b) \binom{a}{2}
 \end{aligned}$$

When $b=0$,

$$\begin{aligned}
 t(n, K_r) &= \binom{n}{2} - (r-1) \binom{\frac{n}{r-1}}{2} \\
 &\approx \left(1 - \frac{1}{r-1}\right) \binom{n}{2} + l.o.t. \\
 &= \left(1 - \frac{1}{r-1}\right) \binom{n}{2} + O(n^2)
 \end{aligned}$$

l.o.t = lower order term

What is $\lim_{n \rightarrow \infty} \frac{t(n, G)}{\binom{n}{2}}$?

Theorem 3 (Erdős-Simonovits-Stone 1946) For any G , $\lim_{n \rightarrow \infty} \frac{t(n, G)}{\binom{n}{2}} = 1 - \frac{1}{\chi(G)-1}$

Theorem 4 (Erdős-Stone) For all integers $r \geq 2$ and $s \geq 1$ and every $\varepsilon > 0$, there exists an integer n_0 such that every graph with $n \geq n_0$ vertices and at least $e(T_{n, r-1}) + \varepsilon n^2$ edges contains $K_r^{(s)}$ as a subgraph.

For any graph H , we define $H^{(s)}$ as follows:

1. $V(H^{(s)}) = \{V(H) \times [s]\}$
2. $E(H^{(s)}) = \{(u, i) \sim (v, j): \text{if } uv \in E(H)\}$

$H^{(s)}$ is called the blow-up graph.

$$K_r^{(s)} = T_{rs, r}$$

Let G be a graph with $\chi(G) = r$ (means G is r -colorable). We can color $V(G)$ so that the endpoints of each edge receive different colors. We can see $G \subseteq K_r^{(s)}$ for some s .

$$t(n, G) \geq \left(1 - \frac{1}{r-1}\right) \binom{n}{2} + O(n^2)$$

Turán graph $T_{n, r-1}$ is G -free and has $\left(1 - \frac{1}{r-1}\right) \binom{n}{2} + O(n^2)$ edges. Because $G \subseteq T_{n, r-1}$, then $\chi(G) \leq r-1$

$$t(n, G) \leq \left(1 - \frac{1}{r-1}\right) \binom{n}{2} + \varepsilon n^2 \text{ for any } \varepsilon > 0.$$

$$1 - \frac{1}{r-1} \leq \frac{t(n, G)}{\binom{n}{2}} \leq 1 - \frac{1}{r-1} + \varepsilon$$

Szemerédi's Regularity Lemma \Rightarrow Blow-up Lemma \Rightarrow Erdős-Stone Theorem
 \Rightarrow Erdős-Simonovits-Stone Theorem

Lemma 2 (Szemerédi's Regularity Lemma) For any $\varepsilon > 0$, and every integer $m \geq 1$, there exists an integer M such that every graph of order at least m admits an ε -regular partition $\{v_0, v_1, \dots, v_k\}$ with $m \leq k \leq M$.

Definition 4 A, B are disjoint sets. The **density** $d(A, B) = \frac{e(A, B)}{|A||B|}$.

Definition 5 (A, B) is called ε -regular if for all $X \subseteq A, Y \subseteq B$ with $|X| \geq \varepsilon|A|$ and $|Y| \geq \varepsilon|B|$,

$$|d(X, Y) - d(A, B)| \leq \varepsilon.$$

(ε is a very small constant $0 < \varepsilon < 1$).

Definition 6 A partition $\{v_0, v_1, \dots, v_n\}$ of $V(G)$ is called ε -regular if

1. $|v_0| \leq \varepsilon|v|$, v_0 -exceptional set
2. $|v_1| = \dots = |v_k|$
3. All but at most εk^2 of the pairs (v_i, v_j) with $1 \leq i < j \leq k$ are ε -regular.

Lemma 3 Suppose (A, B) is ε -regular. Then for any $Y \subseteq B$ with $|Y| \geq \varepsilon|B|$, then all but $\varepsilon|A|$ of the vertices in A have (each) at least $(d - \varepsilon)|Y|$ neighbors in Y .

Proof of Lemma: Let $X = \{v \in A : |N(v) \cap Y| \leq (d - \varepsilon)|Y|\}$. If $|X| \geq \varepsilon|A|$, $e(X, Y) \leq |X|(d - \varepsilon)|Y|$
 $d(X, Y) < d - \varepsilon$ Contradiction. \square

Lemma 4 (Blow-Up) For all $d \in (0, 1]$ and $\Delta \geq 1$, there exists an $\varepsilon_0 > 0$ with the following property: If G is any graph, H is a subgraph with $\Delta(H) \leq \Delta$, $S \in \mathcal{N}$ and R is any regularity graph of G with $\varepsilon \leq \varepsilon_0$, $l \geq \frac{s}{s_0}$ then $H \subseteq R_s \Rightarrow H \subseteq G$. ($\Delta =$ maximum degree)

Proof of Lemma: By regularity lemma,

1. $|v_1| = |v_2| = \dots = |v_k| = l$
2. $|v_0| \leq \varepsilon l$
3. All but εk^2 pairs of (v_i, v_j) are ε -regular.

For vertex $v_i \in V(H)$, we maintain a set of candidates, S_i . $v_i \rightarrow V_{f(i)}$ where $f(i)$ is the index of partition that v_i belongs to.

Initially, $S_i = V_{f(i)}$. For i from 1 to $|H|$ pick $v_i \in S_i$ then update S_j for all $v_j \in N_H(v_i)$.

$$S_j \leftarrow N_G(v_i) \cap S_j$$

$$|S_j^{new}| \geq (d - \varepsilon)|S_j^{old}|$$

For any time, $|S_j| \geq (d - \varepsilon)^\Delta l$. Choose ε such that $(d - \varepsilon)^\Delta \geq (\Delta + 1)\varepsilon$. So, $|S_j| \geq (d - \varepsilon)^\Delta l \geq (\Delta + 1)\varepsilon l$. Since $(\Delta + 1)\varepsilon l \geq \Delta \varepsilon l$, the procedure of selecting v_i will not stop. We construct $H \subseteq G$. \square

Theorem 5 (Erdős-Stone) Suppose G has $e(K_{r-1}^{(s)}) + \lambda n^2$ edges then there exists N for $n > N_\lambda$ G contains $K_r^{(s)}$.

Sketch of proof:

$$K_r^{(s)} \subseteq R^{(s)}$$

$$K_r \subseteq R' \subseteq R$$

We define R' as follows:

1. $v_i v_j \in E(R')$ if (v_i, v_j) are ε -regular and $d(v_i, v_j) \geq \lambda$.

We show R' has enough edges $e(R') > (1 - \frac{1}{r-1})\binom{k}{2}$.